



Evaluation of Blockchain Service Level Agreement (SLA) Using Hyperledger Fabric (HLF)

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Abstract

Different sectors are being revolutionized by distributed ledger technology. According to the 2022 market valuation, Hyperledger is now the second-largest blockchain platform for smart contracts. The creation of numerous apps may be sped up and simplified with smart contracts, but there are certain drawbacks as well. For instance, vulnerability contracts are created intentionally to weaken candor, smart contracts are employed to conduct fraudulent activities, and there are many redundant contracts that squander the efficiency of the system for no real reason. To solve these problems, we provide in this research Service Level Agreement (SLA) for Hyperledger smart contracts. We created Hyperledger smart contracts and focused on how smart contracts and consumers used data. By manually analyzing the transactions, we were able to extract four behavioral characteristics that may be used to differentiate between various contract types. Then, a smart contract is built using these to include 14 fundamental functionalities. We provide a data splitting algorithm for splitting the gathered smart contracts in order to create the experimental dataset. Then, we train and test our dataset using an LSTM network. The comprehensive experimental findings demonstrate that our method can discriminate between various contract types and may be used to identify malicious contracts and detect anomalies with acceptable precision, recall, and F1-score.

1. Introduction

The transactions in digital currencies are recorded using the decentralized, distributed technology known as blockchain (X. Li et al.). Decentralization, permanence, anonymity, and auditability are just a few of the advantages this gives the blockchain (Zheng et al.). As a result, blockchain technology has advanced quickly in recent years, and its use cases have expanded beyond the orig-

inal digital currency distribution transaction to include banking, medical, the Internet of Things, software engineering, and other industries (X. Li et al.). The first successful implementation of blockchain technology was Bitcoin (Nakamoto). In 2015 (Desjardins) and 2016 (X. Li et al.), Bitcoin was hailed as the masterpiece in digital market and commodities, respectively. One of the systems for permissioned blockchains is Hyperledger Fab-

ric (Androulaki, Barger, and Bortnikov). The Linux Foundation created it, and it is open-sourced. Every component serves a unique purpose according to its job. The four primary steps of the transaction flow are endorsement, ordering, validation, and committing. Prior to a network's start, each component has to be tailored and defined according to the network's specifications. To launch their bespoke fabric network, designers need to deal with a multitude of factors surrounding all components. The fabric is made up of a number of different parts, including peer nodes, clients, membership, an ordering service, and Chaincode. Most blockchain networks are built for the purpose of transferring and storing money, and in most cases, this involves producing a "currency" that is used on the network. The word "assets" is used generally in relation to Hyperledger Fabric (D et al.). Cash, real estate, a car, or even an insurance policy are all examples of assets. Its main difference from most other blockchain networks is that it will be used by workers of one or more organizations rather than by individuals. It provides a variety of possibilities rather than a single standard of work for the blockchain network. It may be simpler to construct a network in this fashion so that it may be implemented within the business model of the firm in question.

2. Background

By employing a model, the model and the Calipers benchmark tool, overall performance and delay of HLF version 1 were investigated in (Baliga et al.). The study investigated the effects of several Chaincodes variables and various transactions on transactions performance and latency under micro workloads. By incorporating different numbers of Chaincodes, routes, and neighbors, the authors evaluated the efficacy of Hyperledger Fabric characteristics. The evaluation findings demonstrated how sensitive HLF v1.0 performance is to the Ordered option. Also, the outcomes demonstrated that the HLF v1.0 Compiler was unable to manage a transaction in concurrently using numerous virtual CPUs (vCPUs). That can be viewed as a speed constraint for the system.

Through incorporating additional applications to the platform and nodes scales, Nasir et al. in (Nasir et al.) conducted an experimental efficiency study of two distinct flavors of HLF (v0.6 and v1.0) to exam-

ine the completion time, performance, delay, and adaptability. The findings demonstrated that across all important performance indicators, HLF v1.0 usually beat HLF v0.6.

In (Kuzlu et al.), the authors assessed the HLF v1.4 platform's performance in terms of 'm getting, delay, and scaling with a range of network workloads, including diverse transaction volumes, types, and rates.

The research explained in (Wang) with developed different malicious behaviour patterns examined the effect of malicious conduct on the transactional latency and throughput aspects of HLF performance. The results revealed a considerable decrease in systems efficiency because of assault delay as well as the failure of certain duplicates.

A revolutionary architecture was put forth in Paper (W. Li et al.) to improve the durability of public blockchains. The new design solves the scalability problem of individual Byzantine-fault-tolerant-based (BFT) systems like HLF v0.6 and offers satellites chains to establish a collection of networks. Secure resource transfer across network is possible with suggested design.

A information middle and high the Ethereum Blockchain was created in (Takahashi, Kanai, and Nakazato). A range of process and function, including photos and 3d models, were used to test the study's validity in order to show how well the block chain technology performs overall when used for exchanging sensor information.

2.1. Hyperledger Fabric

Hyper ledger fabric is a blockchain platform with permissions wherein users may see and trust one another. However, it might be set up according to the model of governance developed in order to foster confidence among the users. The orchestration and deployment of distributed ledgers inside coalitions depends on the participation of the collaborating entities. The blockchain is hosted by nodes (or peers), who also execute smart contracts and cooperate to keep apprised of the current state of the record. The common Chaincode can be implemented by all entities or developed privately thanks to the HLF's channel idea. Chain codes may be sent to a set of nodes in a private manner, rendering them inaccessible to other nodes. Only those who subscribe to this very same network have access to the information

and Chaincode. To recognise and verify the nodes, the HLF system must use encrypted resources created by itself. Consequently, this method may be used to authorise certain existing customers. The approved transactions are ordered by the network according to each channel. The approved transactions are ordered by the network according to each channel. Figure 1 displays the comprehensive HLF system design.

2.2. Transaction Flow in Hyperledger Fabric

The transactional method used by HLF is execute-order-validate-and-commit (EOVC). The transactional process of the HFL private ledger is depicted in Figure 2. The stages involved are verification, sorting, and approval. Chaincode spellcasters are Docker (Shahbazi and Byun) container-based transactions. Therefore, separating them from other active chaincodes on the same host as well as the HLF codes would indeed be beneficial. The network nodes maintain a record of transactions as part of the private ledger technology known as HLF. The structure of the data and the status of the data are the two components of the blockchain. The state data describes the nodes' actual situation at any given time, whereas the transaction log records all activities that have been invoked. By running the chaincode, several actions might be performed on the current data status. Transactions are created during processing and added to the log file. The state's statistics may also alter as a result. The LevelDB, a compact library for creating a key-value data store integrated into the HLF node implementation, is used to generate the ledger of transactions. The combinations of session keys make up the status metadata. At the node level, the status database is hot-swappable. A straightforward query for key-value pairs is supported by LevelDB. The CouchDB and NoSQL databases that enable the execution of complicated requests may nevertheless substitute for it. Prior to actually initiating the HLF transaction processing, it is necessary to reveal the credentials of involved stakeholders, associated MSPs, and neighbors. Upon that Orderer system's activation, the network must be started with the appropriate organization's MSPs and nodes joining the channel and initialising the blockchain. Furthermore, the link must have the necessary chaincodes implemented.

2.2.1. Phase 1 — Endorsement Phase:

According to the endorsement rule encoded in the Chaincode, client apps employing the HLF-SDK (Hyper Ledger Fabric Software Development Kit) generate a demand suggestion as well as propose the transaction to endorsing peers. The request is signed cryptographically by the client and sent over similar channels as the blockchain network. The suggested transactions are carried out by the endorsing peers, who also get the predictable output. An endorsement rule must be specified while implementing the Chaincode to determine whose peers (organisations) have had the authority to approve a transaction in the shared smart ledger so that it can be accepted as legitimate and added to the blockchain. To make sure that agreement is reached among all the peers in the process, the HLF takes a number of actions, including establishing endorsement guidelines for ordering services. The sequence of transactions is execute-order-validate-commit (EOVC). Client interactions were first carried out in sandbox environments to ascertain their read-write sets, or the collection of keys that the payment transactions will read from and write to. The transactions are then ordered by an ordering service before being verified and committed to the blockchain. This procedure is carried out by nodes with designated duties. The read-write sets, reply values, and crypto-graphic data that were generated as a consequence of blockchain code processing are all included in the endorsement reply. The endorsing peer delivers a proposed answer to the client after signing the transactional response with their identity using the Chaincode ESCC (Endorsement System Chain Code) system. At this moment, nothing has altered with the ledger. Client application programmes gather recommendations from a variety of endorsing peers and provide them to an ordering phase in accordance with the endorsement rule established by Chaincodes.

2.2.2. Phase 2 - Ordering Phase:

The ordering service audits each transaction in chaincode. They are arranged by network. The transactions that were ordered must then be delivered by the ordering service. A block, which is a transaction transmitted to every peer within the same order by the ordering service, is required by every peer. Any of the established ordering meth-

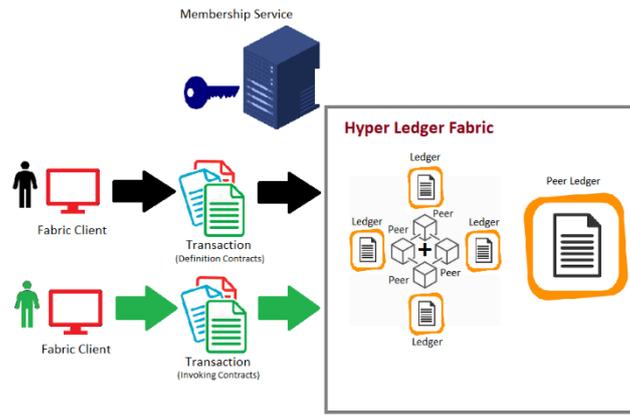


FIGURE 1. HLF network system architecture with its major components

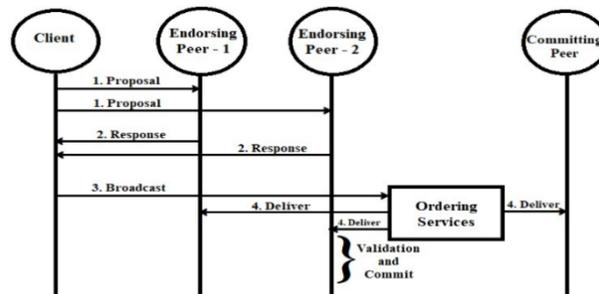


FIGURE 2. Fabric transaction flow.

ods, including Kafka, Solo, or Raft, is used by the ordering service (Shahbazi and Byun). One ordering service node makes up the Solo ordering algorithm. It is employed for system design that runs on a node. The Kafka algorithm is more fully implemented. The Kafka group is created in order to produce and ingest transactions. As a result, it offers crash fault tolerance (CFT) and therefore is advised for a continuous growing list of records in a production setting. The Raft is comparable to the Kafka due to its own commander structure. Its CFT status is a benefit. The raft’s configuration is somewhat simple. In order to see if client nodes can transmit or accept blocks on a specific channel, the ordering service must run access control permission checks on the client nodes. SOLO use is permitted by Hyperledger Fabric version 1.4. It is a message sent from the middle of the gateway, which is simply configured with one node. In a configuration that is better suited for a development stage, no cluster functionality is provided. Nevertheless, the network now has a single failure point. Development teams don’t think about using it in operations.

Kafka would be used in the finished versions of

HLF. High bandwidth and scalability are features of this communication program. It is possible to accept fault tolerance when using a cluster. Multiple communication channels could be tested with different setups because Kafka and SOLO both allow them. A user receives the Orderer to advertise the endorsed transactions to the different peers. The customer receives the responses, which are the endorsements from the endorsing peers. The client is prepared to distribute it to multiple network peers. This step is accomplished by invoking the orderer for the broadcast services. The anchoring peers within that organisation distribute the blocks to the different additional peers after obtaining the blocks containing the transaction. As a result, the person who ordered provides the communication layer of the HLF channel. It is in charge of maintaining the timeline and plays an important role in the consensus protocol.

2.2.3. Phase 3 - Validation Phase:

The transactions must be verified after the blocks has been distributed to all network nodes using the ordering phase provider or chatter protocol. Incorrect transactions are found and rejected through-

out the validation phase. Only legitimate transactions thus are committed as well as maintained in the fabric ledger as well as the world level state. The validation services comprise of two sequential steps: read and write discrepancy checking, often known as Multi-version Concurrency Access Controls (MVCAC), and analysis of endorsement phase using the Validation System Chain Code (VSCC).

2.3. Key Indicators and Tuning Factors

The goal of this work is to examine the efficiency of the modular HLF in a distributed architecture with varied levels of nodes and circumstances in order to determine whether certain factors impact system performance. The research is confined to a deep examination of a few characteristics, while other elements are discussed in broad strokes in order to comprehend and characterise the interrelationships of nodes. As a result, the central emphasis is on examining holistic performance from the vantage point of peers. Concurrently, the Orderer and Gossip effects on the investigation were removed because they were fixed. Figure 3 depicts the overarching test system and its subsystems. A solitary HLF system of one user executing assessment utilities and one anchorage is included in the design.

2.3.1. Performances Parameters:

A report with accurate performance measurements that are relevant to multiple DLT platforms was created by the Hyper-ledger fabric Performance as well as Scalability Workgroup. It was applied inside the experiments and analyses that were described.

2.3.2. Transaction Efficiency:

Smart contracts are deployed, executed, and invoked at varying rates in various public blockchain system. Hence, it is necessary to keep an eye on transaction efficiency. It is calculated as the rate at which the HLF network commits legitimate transactions over a predetermined time frame. The assessment at a single peer is taken into account for the Hyper Ledger Fabric network with one channel. Furthermore, the trials were expanded to include numerous peers in the published framework and analyses (up to 500).

2.3.3. Latency Transaction:

It takes a while for the computer to verify the transactions after it has been transmitted to a connection. The length of time between the period a transaction which performed and the period is confirmed, com-

mitted, and the outcome is made accessible throughout the network is known as the transaction latency. Each transaction counts toward this metric. But in the majority of situations, the experiment offers different statistics on total transactions, including low, high, medium and standard deviations. The transactions verification at a one peer and several peers having various levels of load were examined in the study that was published. Three factors make up the estimated end-to-end transmission delay: the commit, ordering, validation and endorsement phase (Jamil *et al.*).

2.3.4. Scalability:

The capacity of both the deployed HLF network to assist growing the member base is calculated in this study. The scale of the network corresponds to the number of verifying peers taking part in SUT consensus. The limited network is displayed to reflect the overall number of nodes that are currently using the HLF public blockchain network.

2.3.5. Size of the Block:

The three parameters that make up the size of the block—the maximum transaction count per second is calculated, the absolute maximum byte values, and the recommended maximum bytes—present the number of transactions per block seconds. A block of transactions is batch processed. The assessment is further expanded in the study to incorporate batches of numerous blocks value (1 blocks, 10 blocks, 50 blocks and 100 blocks). Additionally, it investigates how various batch sizes affect HLF systems.

3. Environmental Setup

HLF's efficiency and scalability were assessed by implementing several sets of parameters such as transaction sending time, size of the block, size of the network, and NetFlow delivery. The study was performed to evaluate has taken into account a number of measures, including network speed, average sale latency, and available resources. By expanding the connection, scalability was evaluated based on changes in bandwidth and transaction delay. The test findings highlight efficiency bottlenecks, explain the effect of a particular parameter on the HLF public blockchain network, and demonstrate how changes may be made to improve efficiency.

The detailed experimental design used in all of the

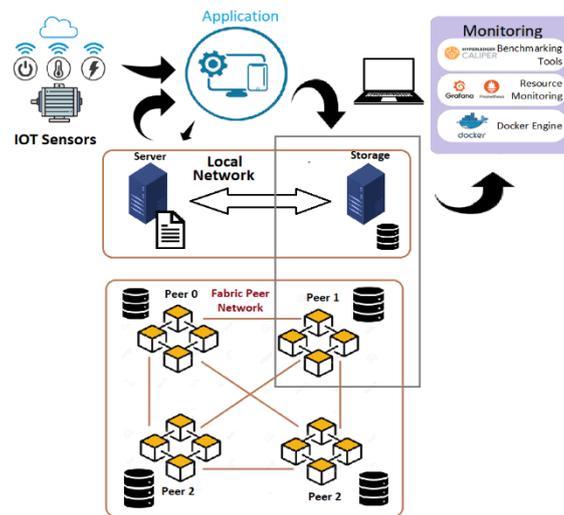


FIGURE 3. HLF-based distributed system model ()

experiments is illustrated in Figure 4. In each case, a public blockchain HLF connection was established, including one organisation and numerous peers. A through-for was used to construct the ordering system, which was executed on a different node. To make the prescribed responsibilities easier, a line of code was put into use on the network.

A permissionless HLF blockchain network was set up in a managed distributed manner. Many Amazon Web services (AWS) EC2 instances were set up as such an underground node-based network in order to get accurate outcomes. Table 1 summarizes the parameters of SUT. Running every instance on its own Virtual Machine (VM). To lessen the impact of connection delays throughout the studies, all VMs were connected to a single subnet. The identical experiment was carried out multiple times with various peer and node values. KV stands for a Key Value that will be communicated to the public blockchain network as a result of a transactions.

IoT gateways in AWS were designed after EC2 instances. The IoT system has adopted different messaging exchanges as chain transactions. This test benchmark system was performed on an AWS EC2 instances with 2vCPUs, 3.0 GHz Xeon Platinum processors, and 4GB RAM. The AWS EC2 machine was running Hyperledger Fabric version v1.4 together with peers, CA, OS, and Ubuntu 18.04 LTS. The effectiveness of the equipment choice (i.e., CPU and RAM) on the bandwidth, delay, and scalability of the developed public blockchain network was examined using that testing environment. VM runs Docker on multiple computer systems to issue

transactions. HLF version 1.4 is a permissioned blockchain network application. To perform multiple chains of code, the VM hosts the HLC (Hyperledger Caliper). The various network peers (from 1 peer to 500 peers) in a scalable environment It deploys Docker to connect Docker swarms to manage the speed of the container. Multiple nodes uses Ubuntu 18.04 LTS OS.

The strength of a Proof of Work (PoW) agreement is evident. because of it is thought to be safest solution for cryptocurrency application because of its pseudo anonymous character. The players in the chain already are familiar with one another in business environments like Distributed systems and telecommunications environments, thus it seemed superfluous in those settings. Consequently, private blockchain blockchains were made for use in businesses. Platforms that employ consensus mechanism that are simpler & utilise fewer resources, such Raft (Ongaro and Ousterhout), which was used in this work.

To reduce this effectiveness, straightforward transactions must be used. Every transaction produces a number that is added to a value according to the system time specified in the blockchain. These variables each produce a key-value pair that also contains a constant rate. Writing the ledger with transactions requires the blockchain. Peers are allowed to make transactions, and they execute in a secure, isolated Docker container. Every action is indeed a write operation that alters the world. The endorsers then independently execute the chaincode, create a transactional report based on the execution

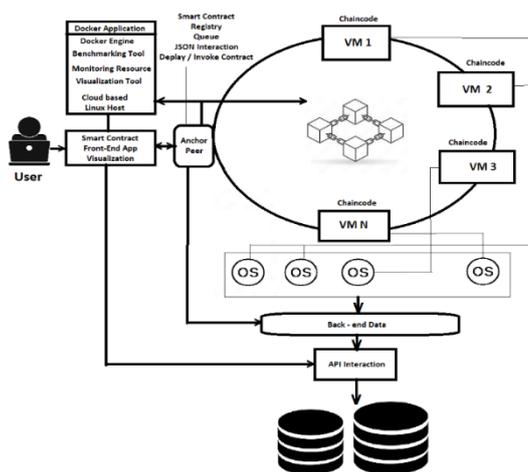


FIGURE 4. Experimental setup for Performance Evolution

TABLE 1. System under test parameters and values

Parameters	Values
Size of Blocks	30 tps
Benchmark out time	1000 ms
Rate of Tx Sending	1-500tps
Number of Blockchain	1,10,50,100 tps
Channels	1 channel
DB State level	Level DB
Transactions	1 KV of 20bytes size
Peers	100v CPU, 3.3 GHZ size, 10 GiB, Moderate network size

outcomes, and verify the reply. The verified transaction request answer is the last thing that the application receives.

4. Summary

To assess the performance and scaling of the Hyperledger Fabric framework in a virtualized environment, many instances with diverse hosts within the network (ranging from 1 to 500 hosts) have been tested. The studies were done in a single-host context with an identical setup. In Hyperledger Fabric version 1, the HLF implemented the concept of different firms. This version 1 was implemented on a multi-host system of virtualized implementations, and the system was measured. Number of operations per unit time, delay, node density, resource utilization, and acceptance policies are examples of key metrics. The outcomes were compared to a solitary host implementation system. It performed better on the majority of the measures. Unfortunately, that approach was not relevant in practice. Simultaneously, the efficiency of multiple-host designs may be improved by deploying numerous organisation

concepts, each with its own systemiser.

The experiments have numerous parameters, such as the peers, size of blocks, and frequency at which transactions are sent, are listed in Table 2. Each test begins with transactions being sent at speeds ranging from 1 to 500 tps.

4.1. Analysis of Single host and Multiple host Transaction

In a single-host configuration, Figure 5 shows the average efficiency of various blocks over varied transaction transmission rates. The average delay for the same transaction sending rates is shown in Figure 6.

In a Multiple-host configuration, Figure 7 shows the average efficiency of various blocks over varied transaction transmission rates and the average efficiency for the same transaction sending rates is shown in Figure 8.

The resulting performance and average delay metrics are shown in Figures 9 and 10. The findings indicate that both situations exhibit a nearly identical delay pattern up to the highest peak. The trans-

TABLE 2. Parameter used for single and multiple transactions of network nodes

Parameters	Values
Peers	1, 5, 10, 20, 30, 40, . . . , 90, 100
Size of the Block	1, 10, 50
Rate of Transactions send	1, 5, 10, 20, 30, 40, . . . , 90, 100, 200, . . . , 400, 500

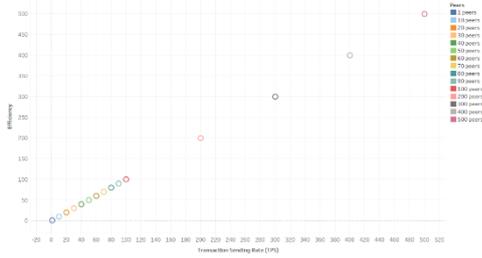


FIGURE 5. Single Host: Transactions sending rate vs Efficiency

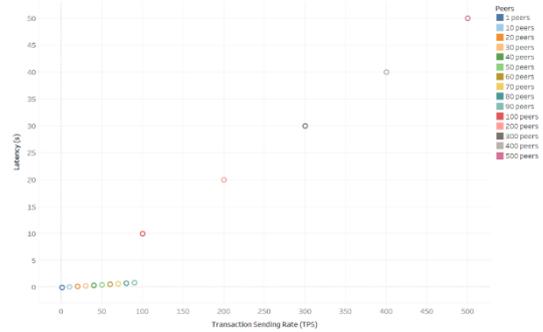


FIGURE 8. Transactions sending rate vs latency for multiple host arrangement.

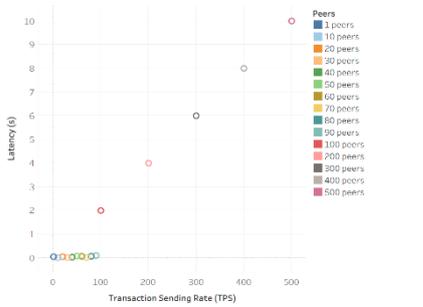


FIGURE 6. Single Host: Transactions sending rate vs latency

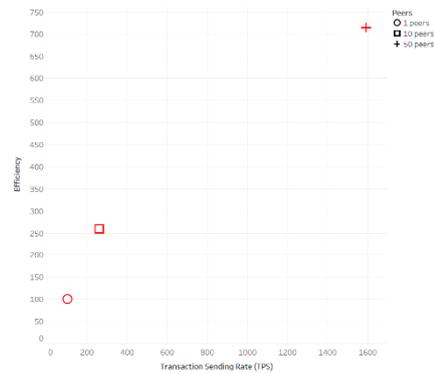


FIGURE 9. Transactions sending rate vs efficiency for different endorsement phase

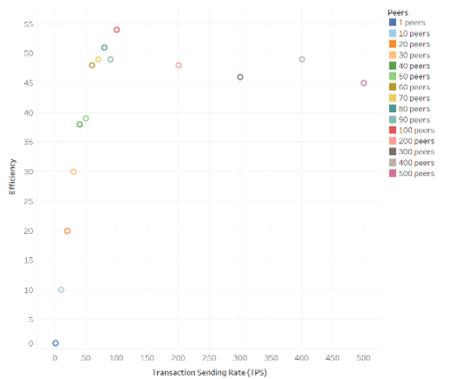


FIGURE 7. Multiple host: Transactions sending rate vs Efficiency

actions delay with a single server host endorsement indicates superior efficiency after the peak point. Despite having varied major position for different configurations, the bandwidth effect increases in each scenario.

Figures 11 and 12 demonstrate that improvement is possible with an improvement in block size. The

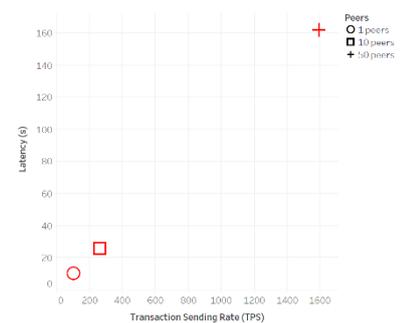


FIGURE 10. Transactions sending rate vs latency for different endorsement phase.

block refers to the number of transactions that can be contained in a block before it is released to the public blockchain network. The Arrange a meeting Batch Size can have a big impact on how fast the system runs. The outcomes demonstrate that what a

small feature size blocks decreases performance.

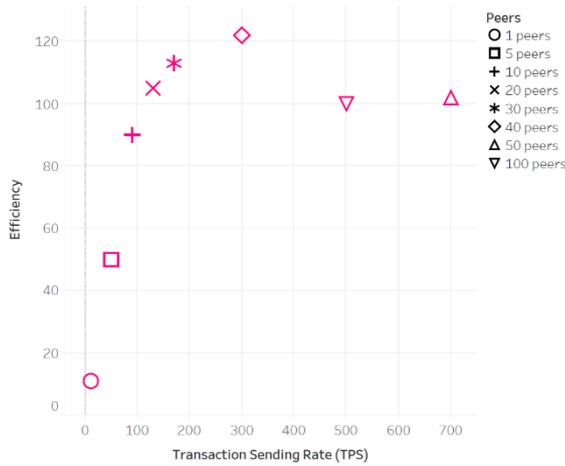


FIGURE 11. Effect of block sizes on system with efficiency.

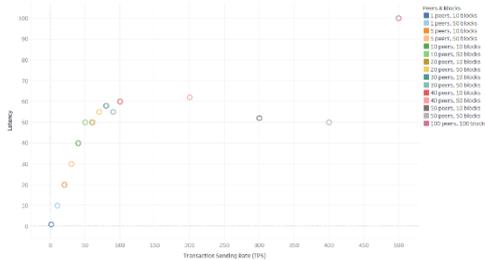


FIGURE 12. Effects of block sizes on transactions with latency

The efficiency and average delays for transactions with various send rates up to 2500 tps are shown in Figures 13 and 14.

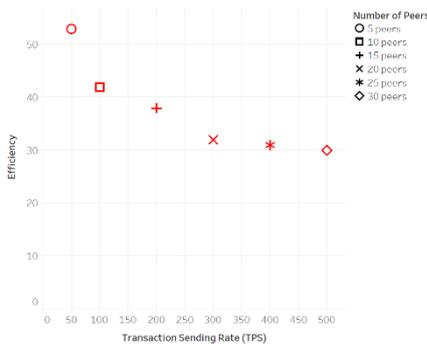


FIGURE 13. Effects of size of the network on the system with efficiency.

The average peer CPU use is shown in Figure 15. Figure 16 illustrates the average disc write utilization and shows the linear increases with small batches as well as the amount of peers. Similar trends of aggregate memory usage by networking neighbours are shown in Figure 17.

Figures 18 & 19 show the Network input and output Traffic of the peers respectively.

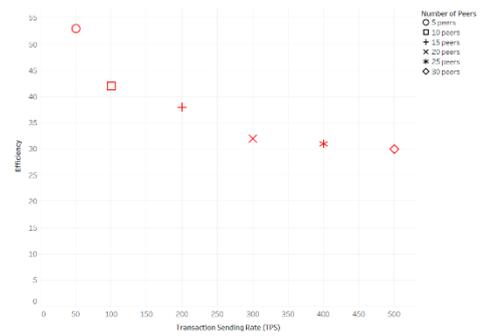


FIGURE 14. Effects of size of the network on transaction with latency

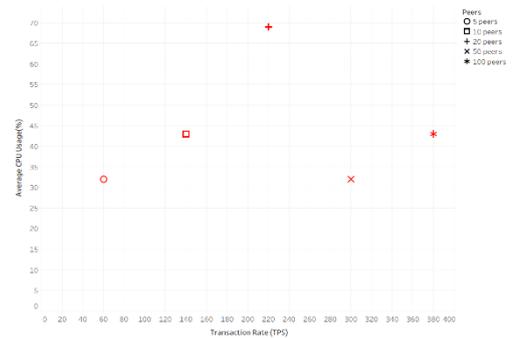


FIGURE 15. Average CPU usage of Peers



FIGURE 16. Average writes disk usage of peers



FIGURE 17. Average memory consumption usage of peers

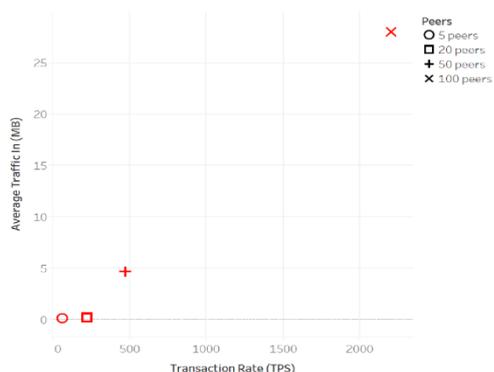


FIGURE 18. Average Network traffic input usage of peers

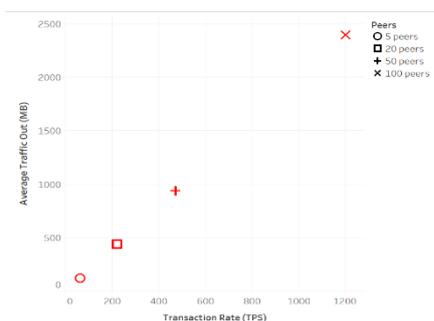


FIGURE 19. Average Network traffic output usage of peers

4.2. Evaluation Metrics

To assess the effectiveness of their systems, we employ the Precision (P), Recall (R), and F1 score (F1) metrics. The term "true positive" (TP) describes how many smart contract forecasts were accurate. False positive (FP) refers to the quantity of incorrectly classifying this kind as an alternative type. Also known as a false negative, this type has been mistakenly classified as a variety of different types. The better it is to discriminate between the various forms of smart contracts, the greater the value of precision, recall, and F1 score.

We have two ways of gathering contract information. One is to synchronise all previously processed transaction information using the client. The alternative is to extract a smart contract's transactions and store them in JSON format using the APIs offered (Ongaro and Ousterhout). We also consult the DApp publication website (Christidis and Devetsikiotis) in accordance with the various uses of smart contracts before classifying the three most prevalent types.

hows the total number of smart contracts in stock market, media and sports. Tables 4-6 showed how

TABLE 3. Number of smart contracts in each different type

Types	Stock Market	Media	Sports
Count	782	338	2930

one type of contract differed from other types, illustrating how the characteristics we glean from smart contract transactions might describe the behavioural patterns of many types of contracts. With respect to these other kinds of contracts, Table 4-6 Precision of sports contracts ranges from 0.902 to 0.955, which indicates that more than 92% of sports of smart contracts can be identified from other kinds by the features. The Recall varied from 0.919 to 0.977, indicating that the more than 92% of contracts with sports characteristics were accurately classified as such. The categorization performance of our models is good, as evidenced by the F1-score, which ranges from 0.916 to 0.982. Our experiment makes use of all 14 of these attributes, and the findings indicate that they can accurately capture the traits of various contract kinds. However, we also discovered that the reliability of the experimental data will be impacted if some smart contracts also have little transactions.

TABLE 4. Result of Stock Market applications with different smart contracts

Type	Precision	Recall	F1-Score
Media	0.932(+/- 0.006)	0.924(+/- 0.009)	0.939(+/- 0.006)
Sports	0.932(+/- 0.019)	0.826(+/- 0.025)	0.876(+/- 0.019)

TABLE 5. Result of Media applications with different smart contracts

Type	Precision	Recall	F1-Score
Media	0.932(+/- 0.006)	0.924(+/- 0.009)	0.939(+/- 0.006)
Sports	0.932(+/- 0.019)	0.826(+/- 0.025)	0.876(+/- 0.019)

Table 7 shows the difference between each one type of contracts which has some differences which varies +/- values is shown for stock market, sports and F1- score.

Focus on F1 score is between 0.701 to 0.835 and overall results are shown using LSTM network

TABLE 6. Result of Sports applications with different smart contracts

Type	Precision	Recall	F1-Score
Stock	0.912(+/-	0.919(+/-	0.916(+/-
Market	0.013)	0.015)	0.013)
Media	0.885(+/-	0.837(+/-	0.860(+/-
	0.025)	0.020)	0.015)

TABLE 7. Results of different types of each single smart contracts.

Type	Precision	Recall	F1-Score
Stock	0.940(+/-	0.858(+/-	0.897(+/-
Market	0.013)	0.016)	0.014)
Sports	0.837(+/-	0.848(+/-	0.876(+/-
	0.013)	0.015)	0.014)
Media	0.888(+/-	0.706(+/-	0.786(+/-
	0.020)	0.026)	0.020)

TABLE 8. Results of different types of smart contracts with LSTM Network

Type	Precision	Recall	F1-Score
Stock	0.805(+/-	0.783(+/-	0.793(+/-
Market	0.052)	0.036)	0.039)
Sports	0.883(+/-	0.698(+/-	0.756(+/-
	0.043)	0.056)	0.039)
Media	0.893(+/-	0.778(+/-	0.826(+/-
	0.047)	0.065)	0.037)

5. Conclusion

This article offered a comprehensive empirical analysis of the extensible Hyperledger Fabric blockchain technology in a decentralized big network with changing peer and payload counts. A scalable framework for accurate and real-time monitoring of HLF systems was presented. The test data demonstrated that the proposed approach was feasible. On the other hand, it revealed that the framework's throughput, delay, and resilience are dependent on system setup, blockchain network architecture, and smart contract complexity procedures. According to the findings, as the volume of transactions and array timeout increases, so does the delay. It is also observed that the number of blocks generated and the volume of transactions for each block had a greater effect on performance, or the number of operations performed per unit time. Due to the number of transactions in a single block, all those transactions are to be verified at the same instant,

resulting in increased performance as block size increases. Furthermore, increasing the array timeout adds delay because each block must wait even though it has completed all transactions.

To optimize effectiveness, experiments with huge network sizes should take into account the optimal values of endorsements each ChainCode to a smaller peer group. For IoT applications with numerous concurrent processes, it might be suggested that higher batch-timeout and block sizes are essential for keeping good throughput. In order to attain low latency, batch timeout as well as block size must be minimal for Iot systems with lots of transactions.

Future work will take into account updating the HLF, assessing the use of legitimate data information, and investigating more test cases, such as examining the effects of maintaining many Orderers just on system's efficiency as a whole. We'll look more closely at other configuration options including having numerous Hyperledger Fabric organisations and having more endorsement peers.

6. Authors' Note

The authors declare that there is no conflict of interest regarding the publication of this article. The authors confirmed that the paper was free of plagiarism.

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